

2-MATCHINGS AND 2-COVERS OF HYPERGRAPHS

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§ 0. Hypergraphs

By a hypergraph \mathcal{H} we mean a finite collection of non-empty finite sets; "collection" means that the same finite set may occur more than once. The elements of a hypergraph are called *edges*; the elements of edges are called *vertices* (this way we exclude "isolated vertices", i.e. vertices not occurring in edges; these vertices would be insignificant in our discussion). The set of vertices of the hypergraph \mathcal{H} will be denoted by $V(\mathcal{H})$.

We need some operations defined on hypergraphs. *Removing an edge* means removing this edge (and, of course, all points which become isolated by this). *Removing a vertex* $x \in V(\mathcal{H})$ means the removal of all edges containing this vertex. *Doubling a vertex* x means the addition of a new vertex x' and of a new edge $E' = E - \{x\} \cup \{x'\}$ for every $E \in \mathcal{H}$, such that $x \in E$. Finally, *cutting off a vertex* x can be carried out if $\{x\} \notin \mathcal{H}$; then it results in the hypergraph $\{E - \{x\}: E \in \mathcal{H}\}$.

A hypergraph obtained from \mathcal{H} by removing edges (points) is called a *partial (induced partial) hypergraph*. The hypergraph obtained by removing the point x is denoted by $\mathcal{H} - x$.

§ 1. Matchings and covers of hypergraphs; results

Let \mathcal{H} be a hypergraph. By a *k-matching of \mathcal{H}* we mean a collection \mathcal{N} of edges of \mathcal{H} (the same edge of \mathcal{H} may occur in \mathcal{N} more than once) such that each point of \mathcal{H} occurs in at most k members of \mathcal{N} . The *size \mathcal{N}* of the *k-matching \mathcal{N}* is the number of its edges.

A *k-cover \mathcal{T}* is a collection of points of \mathcal{H} such that each edge contains at least k of them. The number \mathcal{T} of elements of \mathcal{T} is called the *size of \mathcal{T}* .

By a *fractional matching* we mean a system of weights $n_E \geq 0$ associated with the edges E of \mathcal{H} such that $\sum_{E \ni x} n_E \leq 1$ for each point x . Analogously, a *fractional cover* is a system of weights $t_x \geq 0$ associated with the points x of \mathcal{H} in such a way that $\sum_{x \in E} t_x \geq 1$ for each edge E . The *size of a fractional matching (cover)* is $\sum_E n_E (\sum_x t_x)$.

If \mathcal{N} is a *k-matching* then denoting by N_E the number of times E occurs in \mathcal{N} , $n_E = \frac{1}{k} N_E$ will be a fractional matching. We will denote this fractional matching by $\frac{\mathcal{N}}{k}$. If \mathcal{N}_i is a k_i -matching ($i=1, 2$) then the collection $\mathcal{N}_1 + \mathcal{N}_2$, in which an edge E occurs the number of times it occurs in \mathcal{N}_1 plus the number of times it occurs in \mathcal{N}_2 , is a $(k_1 + k_2)$ -matching. Similar definitions and observations hold for covers.

We denote by $v_k(\mathcal{H})$ the maximum size of k -matching and by $v^*(\mathcal{H})$ the maximum size of a fractional matching. The minimum size of k -covers and fractional covers is denoted by $\tau_k(\mathcal{H})$ and $\tau^*(\mathcal{H})$, respectively. We remove the argument \mathcal{H} if there is only a single hypergraph in consideration. Also we set $\tau_1 = \tau$ and $v_1 = v$.

It is clear from the definitions that

$$(1) \quad v \cong \frac{v_k}{k} \cong v^*$$

and

$$(2) \quad \tau^* \cong \frac{\tau_k}{k} \cong \tau$$

hold for every k . Moreover, the duality theorem of linear programming implies

$$(3) \quad v^* = \tau^*$$

and other well-known results in linear programming imply that there is a k_0 such that $v_{k_0} = k_0 v^*$, $\tau_{k_0} = k_0 \tau^*$. Also, we trivially have $v_{k_1+k_2} \cong v_{k_1} + v_{k_2}$ and $\tau_{k_1+k_2} \cong \tau_{k_1} + \tau_{k_2}$. Hence it is easy to deduce that

$$\lim_{k \rightarrow \infty} \frac{v_k}{k} = \lim_{k \rightarrow \infty} \frac{\tau_k}{k} = v^* = \tau^*$$

and $v_{k_0 \cdot s} = k_0 \cdot s \cdot v^*$, $\tau_{k_0 \cdot s} = k_0 \cdot s \cdot \tau^*$.

Clearly $v(\mathcal{H}) = \tau(\mathcal{H})$ implies that equality holds in (1) and (2) for every k . We can obtain inverses to this observation if we make similar assumption on certain "derived" hypergraphs.

In [4], [5] and [6] the following theorems were shown:¹

THEOREM A. *If $v(\mathcal{H}') = v^*(\mathcal{H}')$ holds for every induced partial hypergraph \mathcal{H}' of \mathcal{H} then $v(\mathcal{H}) = \tau(\mathcal{H})$.*

THEOREM B. *If $\tau(\mathcal{H}') = \tau^*(\mathcal{H}')$ holds for every partial hypergraph \mathcal{H}' of \mathcal{H} then $v(\mathcal{H}) = \tau(\mathcal{H})$.*

THEOREM C. *If $v_2(\mathcal{H}') = 2v(\mathcal{H}')$ holds true for each hypergraph \mathcal{H}' arising from \mathcal{H} by doubling and/or removing vertices then $v(\mathcal{H}) = \tau(\mathcal{H})$.*

Also, the following sharpening of Theorem B is due to Berge.

THEOREM D. *If $\tau_2(\mathcal{H}') = 2\tau(\mathcal{H}')$ holds true for each partial hypergraph \mathcal{H}' of \mathcal{H} then $v(\mathcal{H}) = \tau(\mathcal{H})$.*

The aim of this paper is to discuss the following two theorems analogous to Theorems A—B, and give some applications:

THEOREM 1. *Suppose that $v_2(\mathcal{H}') = 2v^*(\mathcal{H}')$ holds true for each hypergraph \mathcal{H}' arising from \mathcal{H} by doubling and/or removing certain vertices. Then $v_2(\mathcal{H}) = \tau_2(\mathcal{H})$.*

THEOREM 2. *Suppose $\tau_2(\mathcal{H}') = 2\tau^*(\mathcal{H}')$ holds true for each partial hypergraph of \mathcal{H} . Then $v_2(\mathcal{H}) = \tau_2(\mathcal{H})$.*

¹ The first two of these are actually equivalent to certain results of FULKERSON on anti-blocking polyhedra; see [2].